

QuYK: A Universal, Lossless Compression Method for Quantization Matrices

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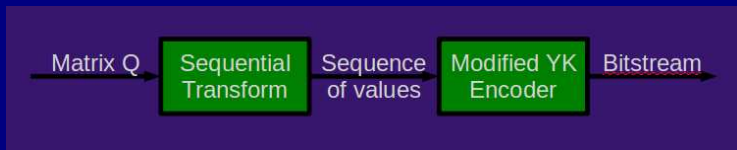
Outline

- 1 Overview and Results
- 2 Grammar-Based Coding
- 3 Conclusions



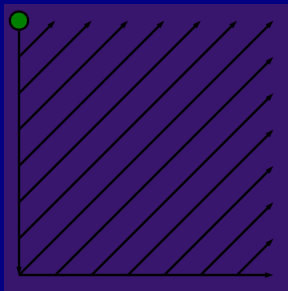
QuYK: A Universal Lossless Compression Method

- Efficient compression of quantization matrices is required to transmit large matrices in the bitstream
- We propose an algorithm using context-free grammars
 - Lossless
 - Universal
 - Provides good compression for the target size
 - Decoding complexity is very low

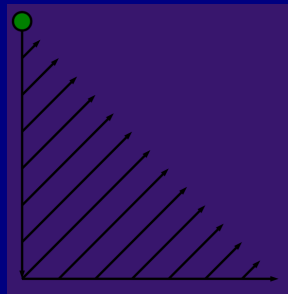


Sequential Transform

- Applies reversible differential coding to the matrix values



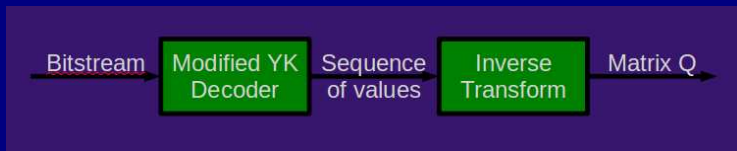
For asymmetric matrices



For symmetric matrices

QuYK Encoding and Decoding

- Encoder is based on a substitutional algorithm using context-free grammars
- Provides lossless reconstruction of the quantization matrix Q
- A special modification makes decoding very simple
 - Version 1 reconstructs the matrix using only string copying
 - Version 2 deploys entropy coding for improved performance



Results

- We demonstrate the performance of our algorithms on the matrices described in JCTVC-D024 [1]

Block size	Number of bits		
	Original	H.264 AVC	QuYK
8x8	512	344	168
16x16	2048	702	432
32x32	8192	1578	664



Existing Lossless Source Codes

- Statistical model-based algorithms and their variants
 - Prediction by Partial Matching
 - Dynamic Markov Coding
 - Context Tree Weighting
 - Process data in symbols
- Dictionary-based algorithms and their variants
 - Lempel-Ziv codes
 - Parse data into phrases



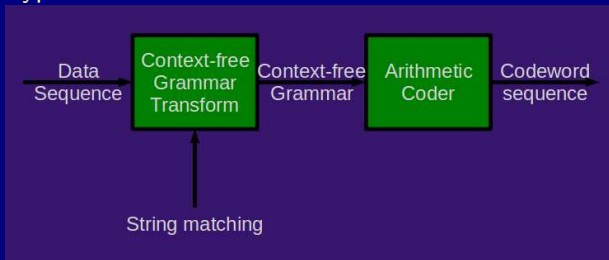
Grammar-Based Coding

- Use admissible grammars for data compression
- First systematic treatment
 - [2] J.C. Kieffer and E.-h. Yang, "Grammar-based codes: A new class of universal lossless source codes," *IEEE Trans. Inf. Theory*, May 2000.
- Context-free grammar-based (CFG-based) codes were proposed in [3]



Context-Free Grammar-Based Codes

- Typical structure of CFG-based code



- Examples
 - Lempel-Ziv algorithms
 - Multi-level pattern matching
 - Context-free Yang-Kieffer algorithms

Context-Free Grammar: An Example

$x = 01023010201$



Context-Free Grammar: An Example

$x = 01023010201$

G_1

$s_0 \rightarrow 01023010201$



Context-Free Grammar: An Example

$x = 01023010201$

G_1

$s_0 \rightarrow 01023010201$

G_2

$s_0 \rightarrow s_1 3 s_1 s_2$

$s_1 \rightarrow s_2 0 2$

$s_2 \rightarrow 0 1$



Context-Free Grammar: An Example

$x = 01023010201$

G_1

$s_0 \rightarrow 01023010201$

CFG size: 11

G_2

$s_0 \rightarrow s_1 3 s_1 s_2$

$s_1 \rightarrow s_2 0 2$

$s_2 \rightarrow 0 1$

CFG size: 9



Context-Free Grammar Transform

$s_0 \rightarrow \underline{01023010201}$



Context-Free Grammar Transform

$$s_0 \rightarrow \underline{01023010201} \implies \begin{array}{lcl} s_0 & \rightarrow & s_1 3 s_1 \underline{01} \\ s_1 & \rightarrow & \underline{0102} \end{array}$$



Context-Free Grammar Transform

$$s_0 \rightarrow \underline{01023010201} \implies \begin{array}{l} s_0 \rightarrow s_1 3 s_1 \underline{01} \\ s_1 \rightarrow \underline{0102} \end{array} \implies \begin{array}{l} s_0 \rightarrow s_1 3 s_1 s_2 \\ s_1 \rightarrow s_2 02 \\ s_2 \rightarrow 01 \end{array}$$



Context-Free Grammar Transform

$$s_0 \rightarrow \underline{01023010201} \implies \begin{array}{l} s_0 \rightarrow s_1 3 s_1 \underline{01} \\ s_1 \rightarrow \underline{0102} \end{array} \implies \begin{array}{l} s_0 \rightarrow s_1 3 s_1 s_2 \\ s_1 \rightarrow s_2 02 \\ s_2 \rightarrow 01 \end{array}$$

- String matching
- The resulting CFG is a compact representation of the original sequence

$$|G| = O\left(\frac{|x|}{\log |x|}\right)$$



Performance for Sources with Finite Alphabets: Redundancy Result

- **Result 1** A CFG-based code's worst case redundancy against the k -context empirical entropy among all length n individual sequence from a **finite** alphabet is upper bounded by

$$\max_{x \in A^n} \{r^\phi(x) - r_k^*(x)\} = O\left(\frac{\log \log n}{\log n}\right)$$

$r^\phi(x)$ the entropy of CFG-based codes followed by arithmetic coding

$r_k^*(x)$ the k -context empirical entropy of x , which represents the smallest compression rate in bits per symbol among all arithmetic coding algorithms with k contexts



Performance for Sources with Finite Alphabets: Universality

- **Result 2** Let A denote a **finite** alphabet. A CFG-based code is **universal** with respect to the class of all stationary, ergodic sources with alphabet A .

$X = \{X_i\}_{i=1}^{\infty}$: a stationary, ergodic source from A

$$r^{\phi}(X_1 X_2 \dots X_n) \rightarrow H(X)$$

with probability 1 as n approaches infinity.



Remarks

- CFG-based codes have
 - Good theoretical performance
 - Good performance in practice (context-free Yang-Kieffer algorithms)
 - Very low computational complexity at the decoder's side



Grammar Transform for QuYK

- Generate the final irreducible grammar for the sequence
- Reorder the production rules for instantaneous access at decoding
 - Rules are written into bitstream in a certain order
 - Any rule may refer only to rules preceding it
- Write the grammar into the bitstream

Clear coding Use approx. $\log|A|$ bits for a joint
symbol-variable alphabet A

Arithmetic coding Improved compression



Conclusions

- Efficient compression of quantization matrices is required to transmit those matrices in the bitstream
- A universal algorithm using context-free grammars provides good performance, at very low decoding complexity
- We propose to study further in the quantization ad-hoc group the compression of quantization matrices, and to define test conditions for their compression



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